

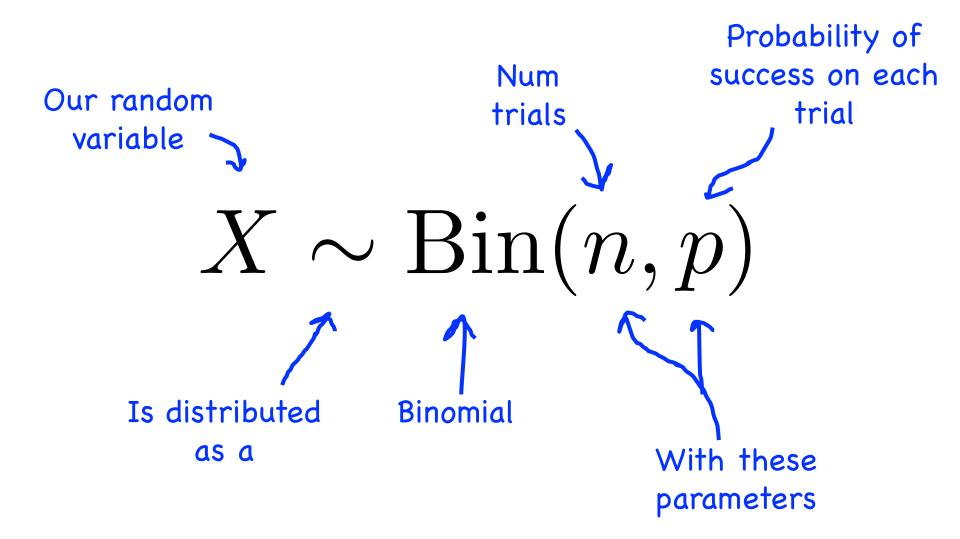
# **HURRICANE MICHAEL** 11:00 AM CDT LAT: 29.6°N LON: 85.8°W 35 MI SW OF MEXICO BEACH FLORIDA Birmingham WINDS: 150 MPH Macon PRESSURE: 923 mb **MOVING: NNE at 14 MP** Panama City Probability for Extreme Weather?

# Review

## **Binomial Random Variable**

- Consider n independent trials of an experiment with success probability p.
  - X is number of successes in n trials
  - X is a Binomial Random Variable:

- Examples
  - # of heads in n coin flips
  - # of 1's in randomly generated length n bit string
  - # of disk drives crashed in 1000 computer cluster
    - Assuming disks crash independently



#### If X is a binomial with parameters n and p

Probability Mass Function for a Binomial

$$P(X=k) = \binom{n}{k} p^k (1-p)^{n-k}$$

Probability that our variable takes on the value k

## Bernoulli vs Binomial

$$X \sim \mathrm{Bern}(p)$$

$$X \in \{0, 1\}$$

Bernoulli is a type of RV that can take on two values, 1 (for success) with probability p and 0 (for failure) with probability (1- p)

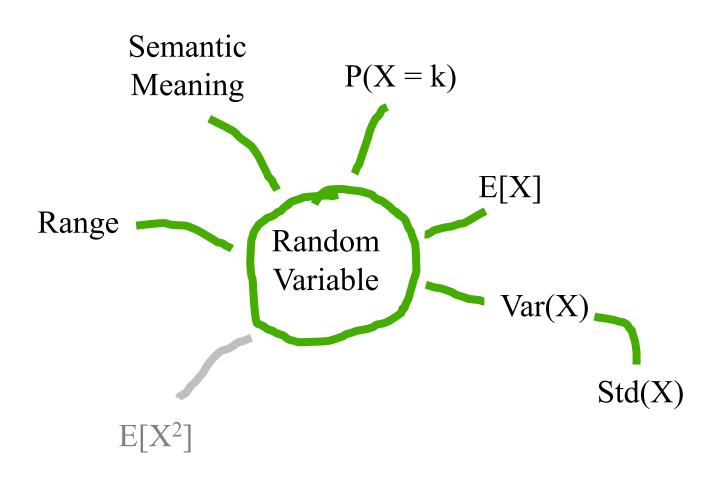
$$Y \sim \text{Bin}(n, p)$$

$$Y = \sum_{i=1}^{n} X_i$$

s.t. 
$$X_i \sim \text{Bern}(p)$$

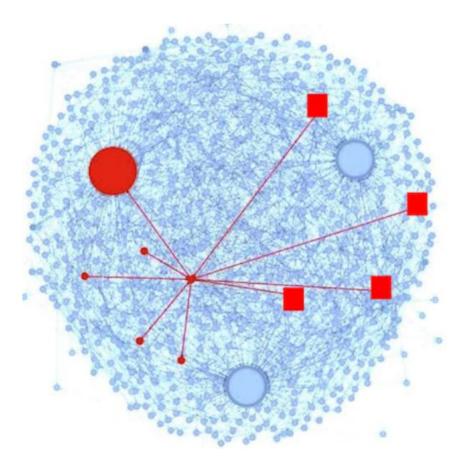
Binomial is the sum of *n*Bernoullis

# **Fundamental Properties**



# Is Peer Grading Accurate Enough?

#### Looking ahead

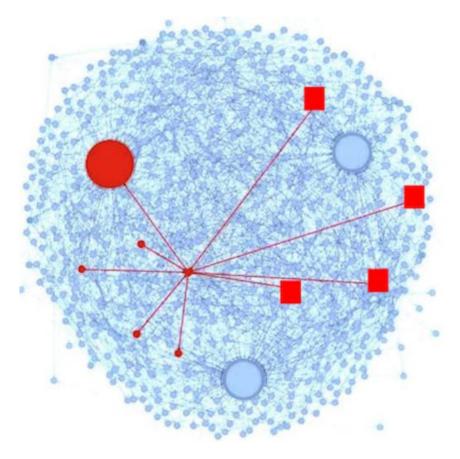


Peer Grading on Coursera HCI.

31,067 peer grades for 3,607 students.

# Is Peer Grading Accurate Enough?

#### Looking ahead



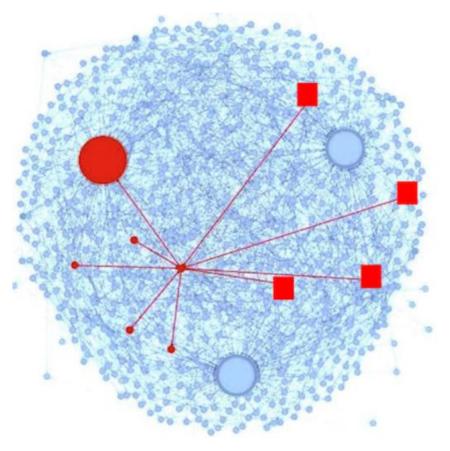
- 1. Defined random variables for:
  - True grade (s<sub>i</sub>) for assignment i
  - Observed  $(z_i^j)$  score for assign i
  - Bias  $(b_i)$  for each grader j
  - Variance  $(r_i)$  for each grader j
- 2. Designed a probabilistic model that defined the distributions for all random variables

variables 
$$s_i \sim \text{Bin}(\text{points}, \theta)$$

$$z_i^j \sim \mathcal{N}(\mu = s_i + b_j, \sigma = \sqrt{r_j})$$

# Is Peer Grading Accurate Enough?

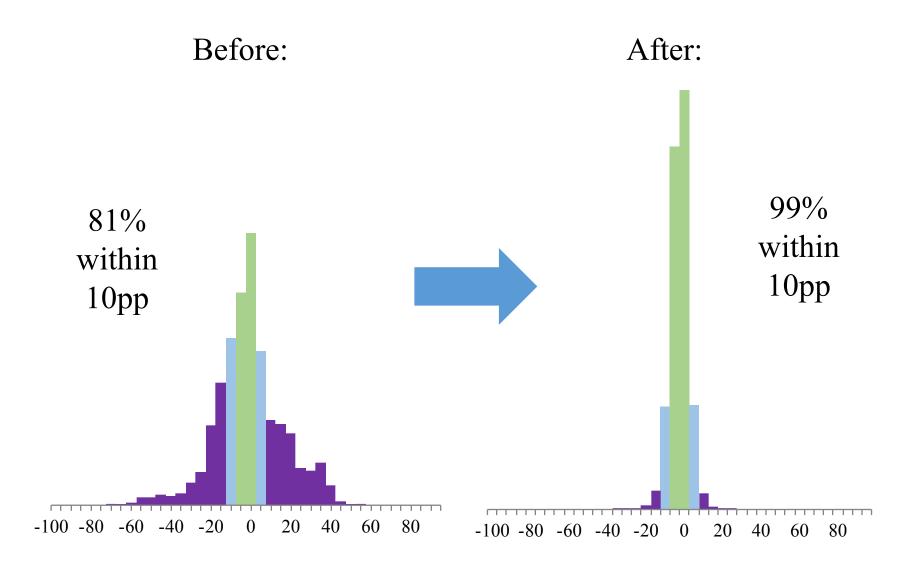
#### Looking ahead



- 1. Defined random variables for:
  - True grade (s<sub>i</sub>) for assignment i
  - Observed  $(z_i^j)$  score for assign i
  - Bias (b<sub>i</sub>) for each grader j
  - Variance  $(r_i)$  for each grader j
- 2. Designed a probabilistic model that defined the distributions for all random variables
- **3.** Found the variable assignments that maximized the probability of our observed data

Inference or Machine Learning

# Yes, With Probabilistic Modelling



Tuned Models of Peer Assessment. C Piech, J Huang, A Ng, D Koller

# **End Review**

## Good to know

#### **Natural Exponent def:**

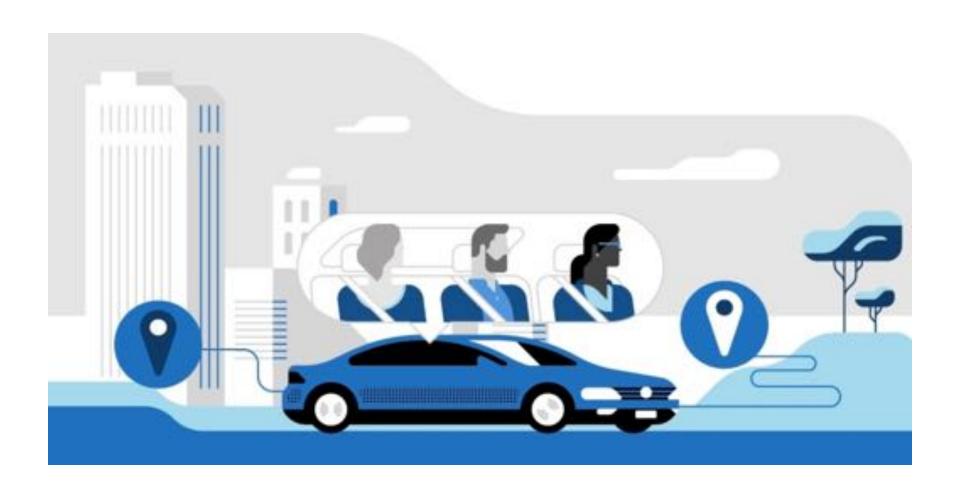
$$\lim_{n \to \infty} \left( 1 - \frac{\lambda}{n} \right)^n = e^{-\lambda}$$

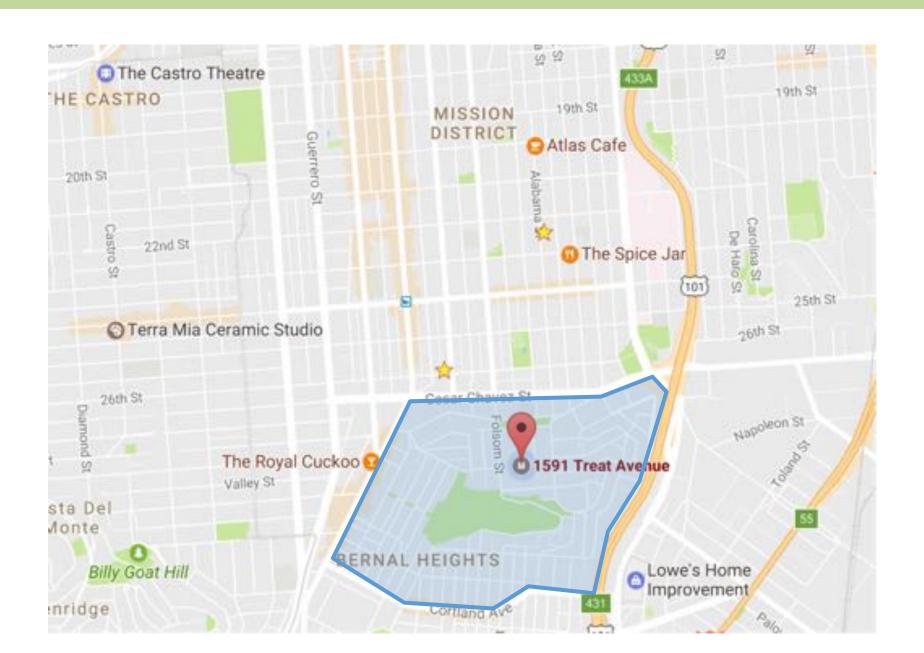
Jacob Bernoulli

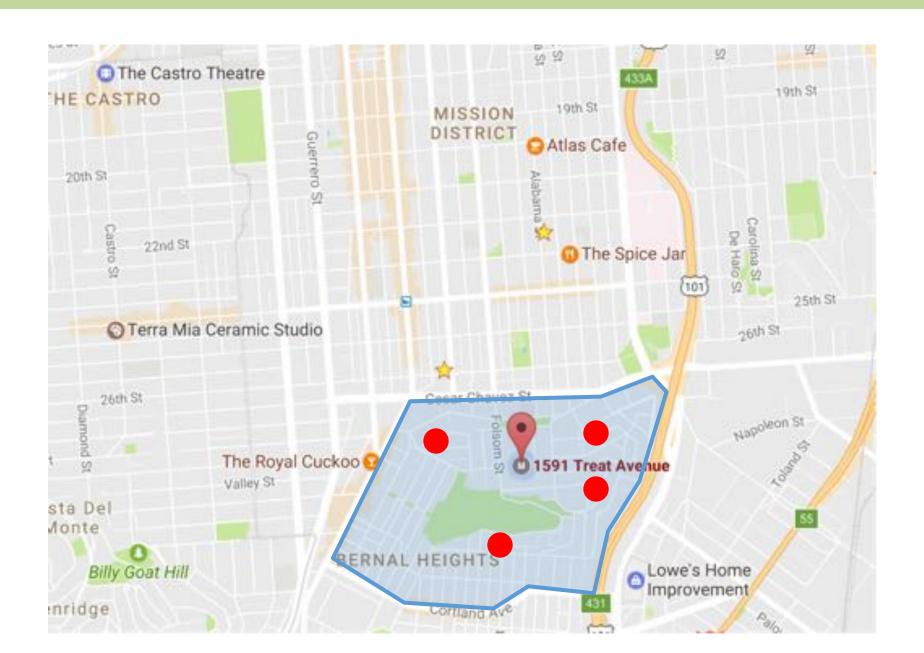


https://en.wikipedia.org/wiki/E (mathematical constant)

# Algorithmic Ride Sharing



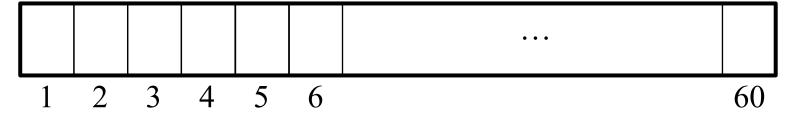






On average  $\lambda = 5$  requests per minute

We can break the next minute down into seconds



On average  $\lambda = 5$  requests per minute

We can break the next minute down into seconds



At each second either get a request or you don't.

On average  $\lambda = 5$  requests per minute

We can break the next minute down into seconds



At each second either get a request or you don't. Let X =Number of requests in the minute

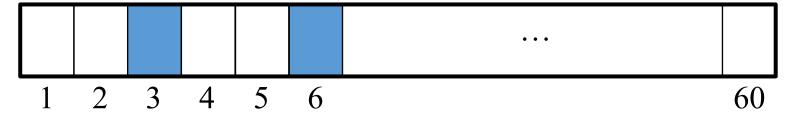
$$X \sim \text{Bin}(n = 60, p = 5/60)$$

$$P(X = k) = \binom{n}{k} (p)^k (1 - p)^{n - k}$$

$$P(X = 3) = \binom{60}{3} (5/60)^3 (1 - 5/60)^{57}$$

On average  $\lambda = 5$  requests per minute

We can break the next minute down into seconds



At each second either get a request or you don't. Let X =Number of requests in the minute

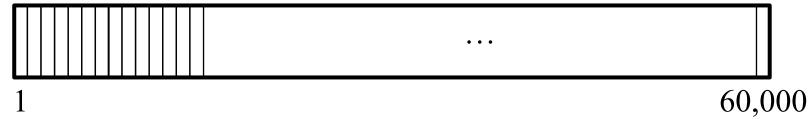
$$X \sim \text{Bin}(n = 60, p = 5/60)$$

$$P(X = k) = \binom{n}{k} (p)^k (1 - p)^{n-k}$$

But what if there are two requests in the same second?

On average  $\lambda = 5$  requests per minute

We can break that next minute down into *milli*-seconds

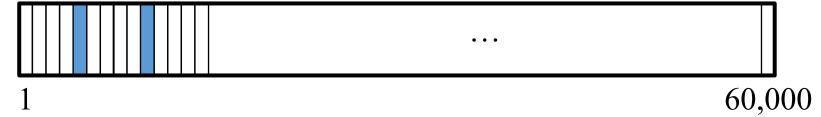


At each *milli*-second either get a request or you don't. Let X = Number of requests in the minute

But what if there are two requests in the same second?

On average  $\lambda = 5$  requests per minute

We can break that next minute down into *milli*-seconds



At each *milli*-second either get a request or you don't. Let X = Number of requests in the minute

$$X \sim \text{Bin}(n = 60000, p = \lambda/n)$$

$$P(X = k) = \binom{n}{k} (\lambda/n)^k (1 - \lambda/n)^{n-k}$$

Can we do any better than milli-seconds?

## Binomial in the Limit

On average  $\lambda = 5$  requests per minute

We can break that minute down into infinitely small buckets

#### OMG so small

1

 $\infty$ 

Let X = Number of requests in the minute

$$X \sim \text{Bin}(n, p = \lambda/n)$$

$$P(X = k) = \lim_{n \to \infty} \binom{n}{k} (\lambda/n)^k (1 - \lambda/n)^{n-k}$$

Who wants to see some cool math?

## Binomial in the Limit

$$P(X = k) = \lim_{n \to \infty} \binom{n}{k} (\lambda/n)^k (1 - \lambda/n)^{n-k}$$

$$= \lim_{n \to \infty} \frac{n!}{(n-k)!k!} \cdot \frac{\lambda^k}{n^k} \cdot \frac{(1-\lambda/n)^n}{(1-\lambda/n)^k}$$
 By expanding each term

$$= \lim_{n \to \infty} \frac{n!}{(n-k)!k!} \cdot \frac{\lambda^k}{n^k} \cdot \frac{e^{-\lambda}}{1}$$

By definition of natural exp

$$= \lim_{n \to \infty} \frac{n!}{(n-k)!n^k} \cdot \frac{\lambda^k}{k!} \cdot \frac{e^{-\lambda}}{1}$$

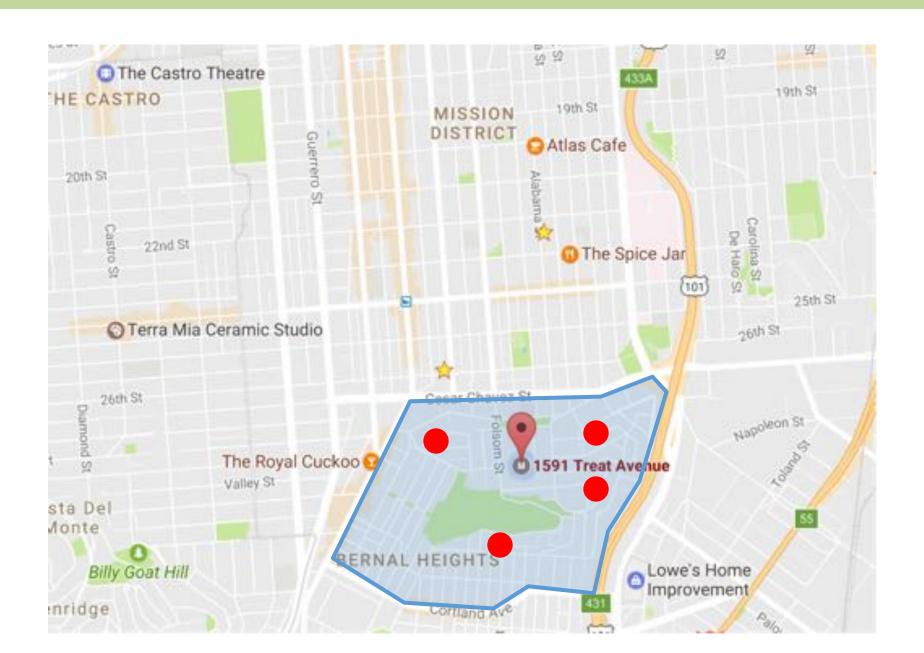
Rearranging terms

$$= \lim_{n \to \infty} \frac{n^k}{n^k} \cdot \frac{\lambda^k}{k!} \cdot \frac{e^{-\lambda}}{1}$$

Limit analysis

$$= \frac{\lambda^k e^{-\lambda}}{k!}$$

Simplifying



## Simeon-Denis Poisson

 Simeon-Denis Poisson (1781-1840) was a prolific French mathematician





- Published his first paper at 18, became professor at 21, and published over 300 papers in his life
  - He reportedly said "Life is good for only two things, discovering mathematics and teaching mathematics."
- I'm going with French Martin Freeman

## Poisson Random Variable

 X is a <u>Poisson</u> Random Variable: the number of occurrences in a fixed interval of time.

$$X \sim \text{Poi}(\lambda)$$

- λ is the "rate"
- X takes on values 0, 1, 2...
- has distribution (PMF):

$$P(X = k) = e^{-\lambda} \frac{\lambda^k}{k!}$$

### **Poisson Process**

- Consider events that occur over time
  - Earthquakes, radioactive decay, hits to web server, etc.
  - Have time interval for events (1 year, 1 sec, whatever...)
  - Events arrive at rate: λ events per interval of time
- Split time interval into n → ∞ sub-intervals
  - Assume at most one event per sub-interval
  - Event occurrences in sub-intervals are independent
  - With many sub-intervals, probability of event occurring in any given sub-interval is small
- N(t) = # events in original time interval ~ Poi(λ)



# Poisson is great when you have a rate!





Poisson is great when you have a rate and you care about # of occurrences!

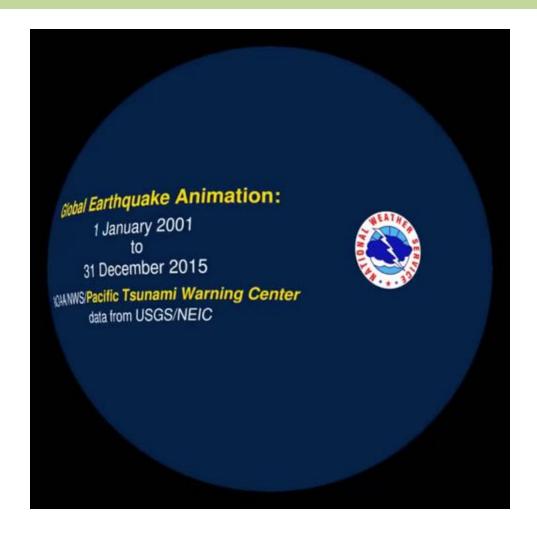




Make sure that the time unit for "rate" and match the probability question



# Earthquakes



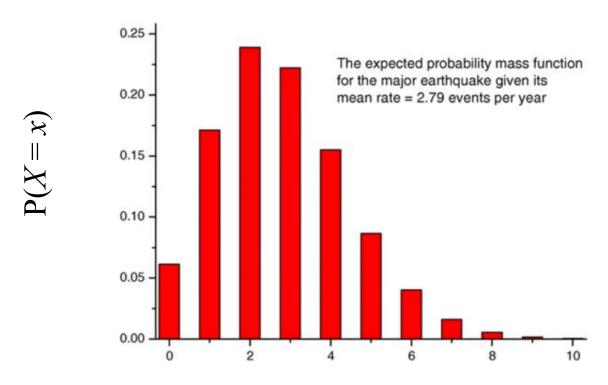
Average of 2.79 major earthquakes per year. What is the probability of 3 major earthquakes next year?



## Earthquake Probability Mass Function

Let X = number of earthquakes next year

$$X \sim \text{Poi}(2.79)$$



Number of earthquakes (x)

$$P(X=3) = \frac{\lambda^k e^{-\lambda}}{k!} = \frac{2.79^3 e^{-2.79}}{3!} \approx 0.23$$



## Bulletin of the Seismological Society of America

Vol. 64 October 1974 No. 5

#### IS THE SEQUENCE OF EARTHQUAKES IN SOUTHERN CALIFORNIA, WITH AFTERSHOCKS REMOVED, POISSONIAN?

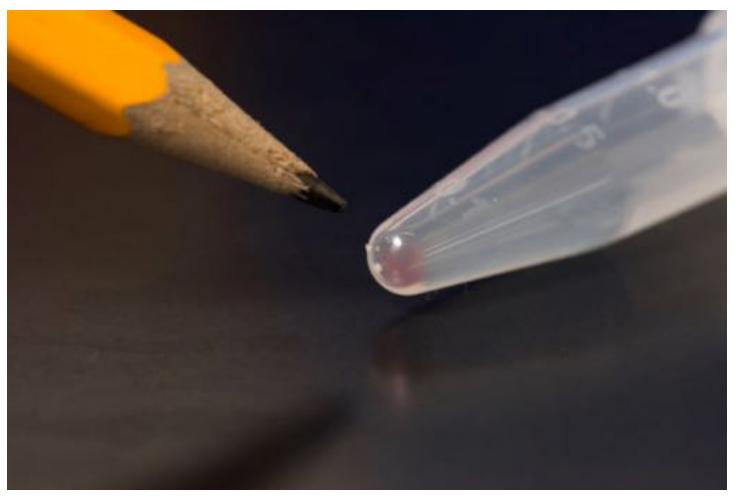
By J. K. GARDNER and L. KNOPOFF

ABSTRACT

Yes.

Poisson can approximate a Binomial!

### Storing Data on DNA



All the movies, images, emails and other digital data from more than 600 smartphones (10,000 gigabytes) can be stored in the faint pink smear of DNA at the end of this test tube.

### Storing Data on DNA

- Will the DNA storage become corrupt?
  - In DNA (and real networks) store large strings
  - Length  $n \approx 10^4$
  - Probability of corruption of each base pair is very small  $p \approx 10^{-6}$
  - X ~ Bin(10<sup>4</sup>, 10<sup>-6</sup>) is unwieldy to compute

- Extreme n and p values arise in many cases
  - # bit errors in steam sent over a network
  - # of servers crashes in a day in giant data center

### Storing Data on DNA

- Will the DNA storage become corrupt?
  - In DNA (and real networks) store large strings
  - Length  $n \approx 10^4$
  - Probability of corruption of each base pair is very small  $p \approx 10^{-6}$
  - $X \sim Poi(\lambda = 10^4 * 10^{-6} = 0.01)$

$$P(X = k) = e^{-\lambda} \frac{\lambda^k}{k!}$$

$$P(X = 0) = e^{-\lambda} \frac{1}{0!}$$

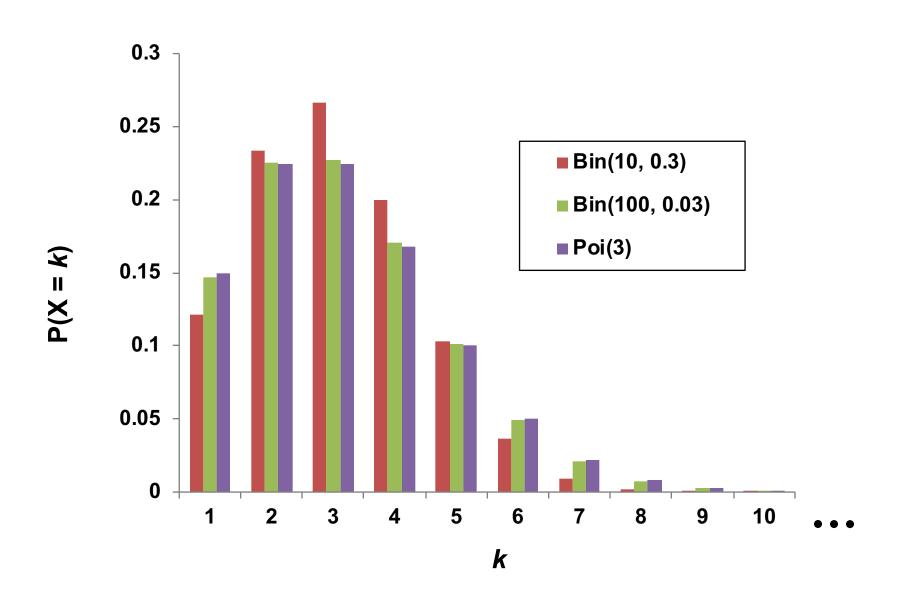
$$= e^{-0.01} \approx 0.99$$

#### Poisson is Binomial in the Limit

- Poisson approximates Binomial where n is large, p is small, and  $\lambda = np$  is "moderate"
- Different interpretations of "moderate"
  - n > 20 and p < 0.05
  - n > 100 and p < 0.1

• Really, Poisson is Binomial as  $n \to \infty$  and  $p \to 0$ , where  $np = \lambda$ 

## Bin(10,0.3) vs Bin(100,0.03) vs Poi(3)





Poisson can be used to approximate a Binomial where n is large and p is small.



#### Tender (Central) Moments with Poisson

- Recall:  $Y \sim Bin(n, p)$ 
  - E[Y] = np
  - Var(Y) = np(1 p)
- $X \sim Poi(\lambda)$  where  $\lambda = np$   $(n \rightarrow \infty \text{ and } p \rightarrow 0)$ 
  - $E[X] = np = \lambda$
  - $Var(X) = np(1-p) = \lambda(1-0) = \lambda$
  - Yes, expectation and variance of Poisson are same
    - It brings a tear to my eye...

#### A Real License Plate Seen at Stanford



No, it's not mine... but I kind of wish it was.

#### Poisson is Chill

- Poisson can still provide a good approximation even when assumptions are "mildly" violated
- "Poisson Paradigm"
- Can apply Poisson approximation when...
  - "Successes" in trials are not entirely independent
    - Example: # entries in each bucket in large hash table
  - Probability of "Success" in each trial varies (slightly)
    - Small relative change in a very small p
    - Example: average # requests to web server/sec. may fluctuate slightly due to load on network

#### **Web Server Load**

- Consider requests to a web server in 1 second
  - In past, server load averages 2 hits/second
  - X = # hits server receives in a second
  - What is P(X < 5)?
- Solution

$$X \sim \text{Poi}(\lambda = 2)$$

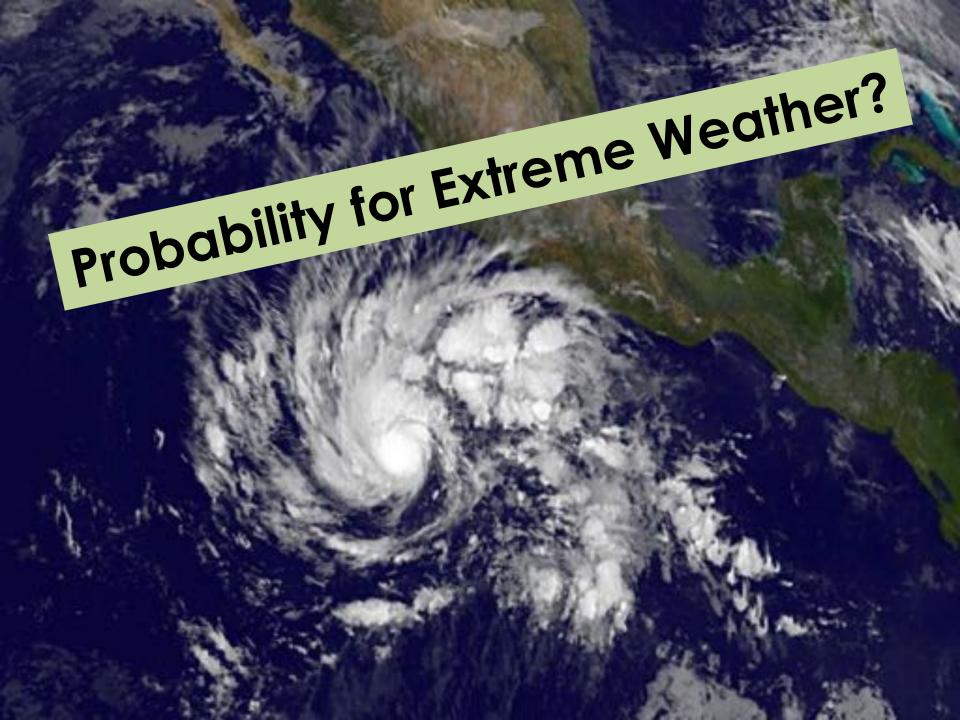
$$P(X < 5) = \sum_{i=0}^{4} P(X = i)$$

$$=\sum_{i=0}^{4}e^{-\lambda}\frac{\lambda^{i}}{i!}$$

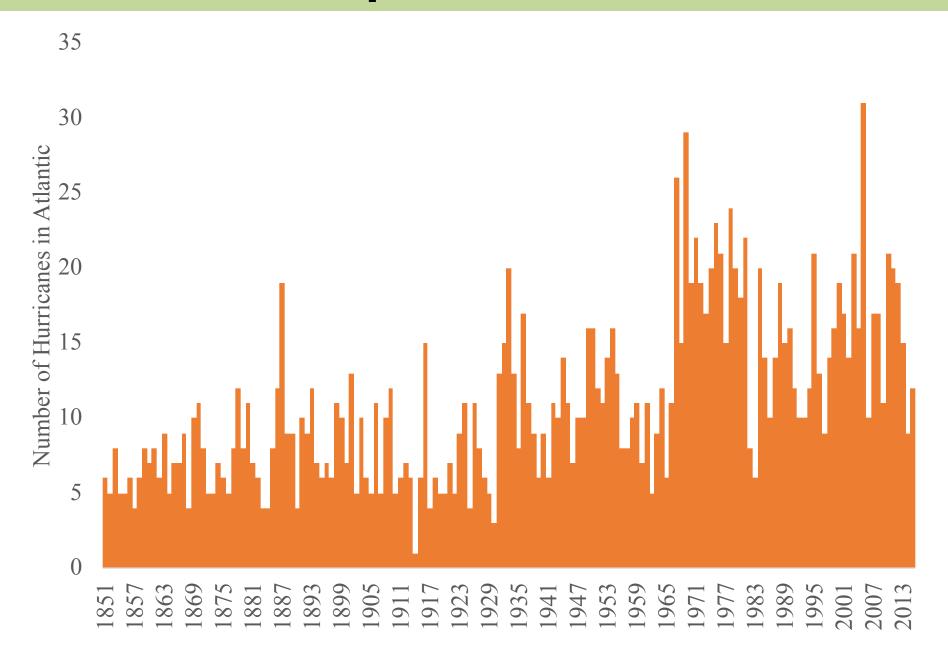
$$=\sum_{i=0}^{4}e^{-2}\frac{2^{i}}{i!}\approx 0.95$$

Since X is Poisson

Since 
$$\lambda = 2$$

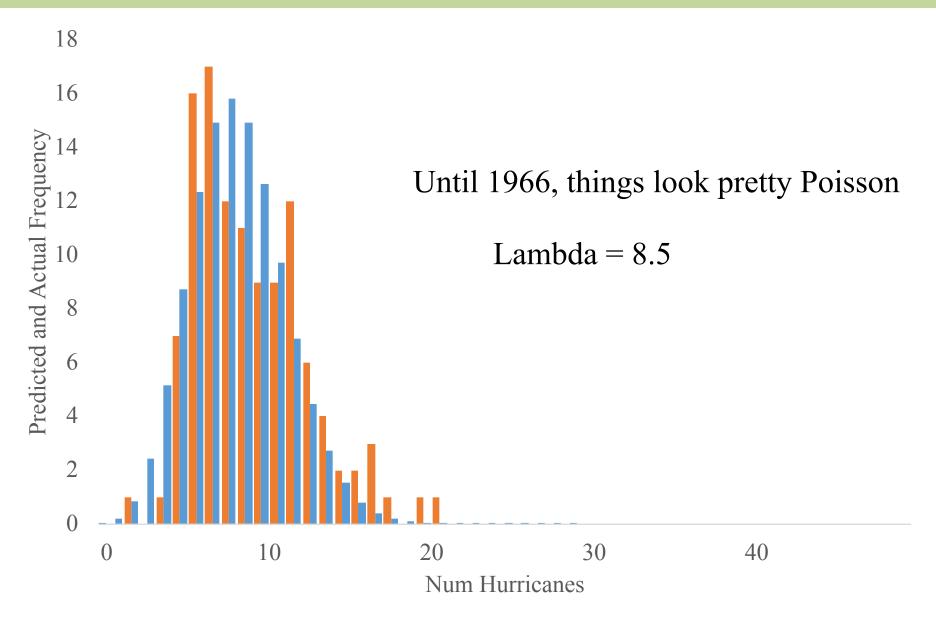


#### Hurricanes per Year since 1851



# To the code!

# Historically ~ Poisson(8.5)



### Improbability Drive

- What is the probability of over 15 hurricanes in a season given that the distribution doesn't change?
  - Let X = # hurricanes in a year. X ~ Poi(8.5)
- Solution:

$$P(X>15) = 1 - P(X \le 15) \qquad \text{of a Poisson.}$$
 Your favorite programming language has a function for it 
$$= 1 - 0.98$$
 
$$= 0.02$$

This is the pmf of a Poisson. Your favorite function for it

# Twice since 1966 there have been years with over 30 hurricanes

### Improbability Drive

- What is the probability of over 30 hurricanes in a season given that the distribution doesn't change?
  - Let X = # hurricanes in a year. X ~ Poi(8.5)
- Solution:

$$P(X > 30) = 1 - P(X \le 30)$$

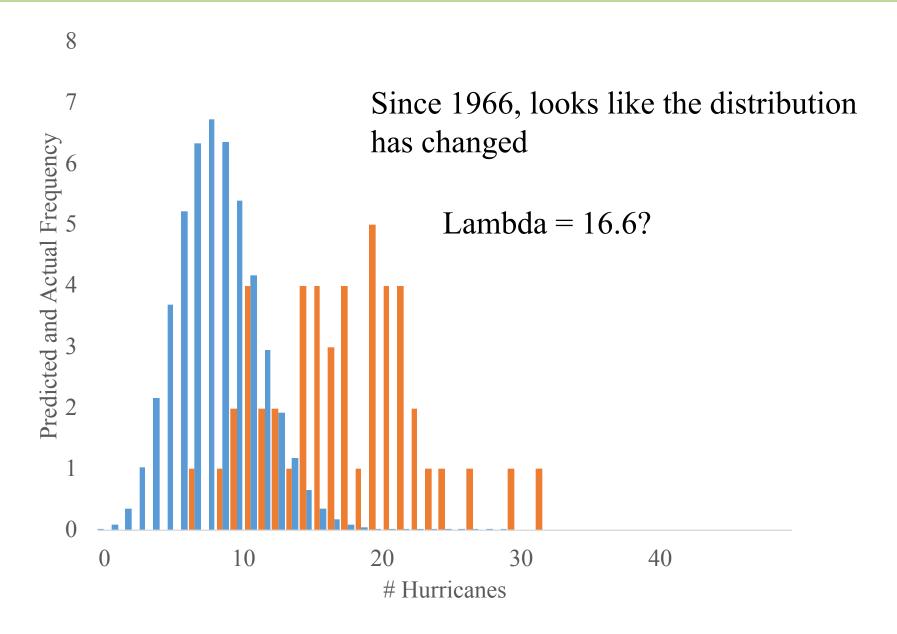
$$= 1 - \sum_{i=0}^{30} P(X = i)$$

$$= 1 - 0.999999997823$$

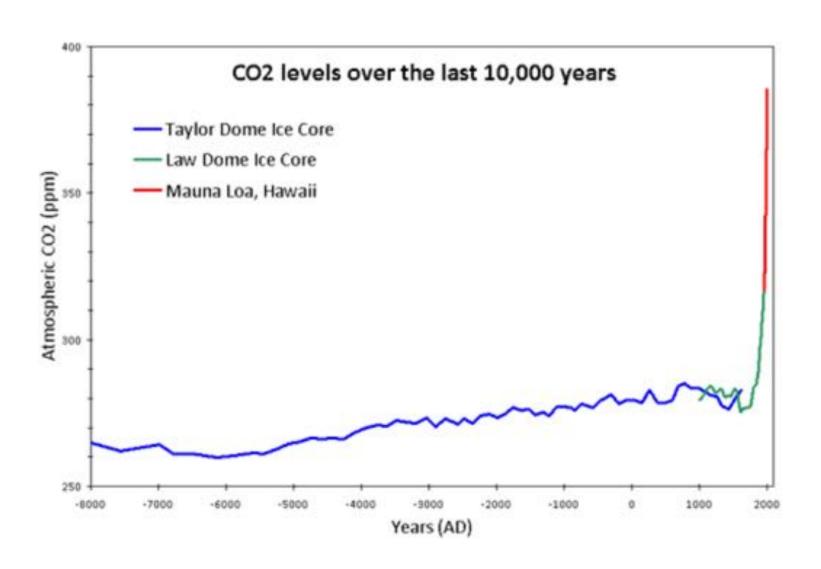
$$= 2.2e - 09$$

This is the pdf of a Poisson. Your favorite programming language has a function for it

### The Distribution has Changed



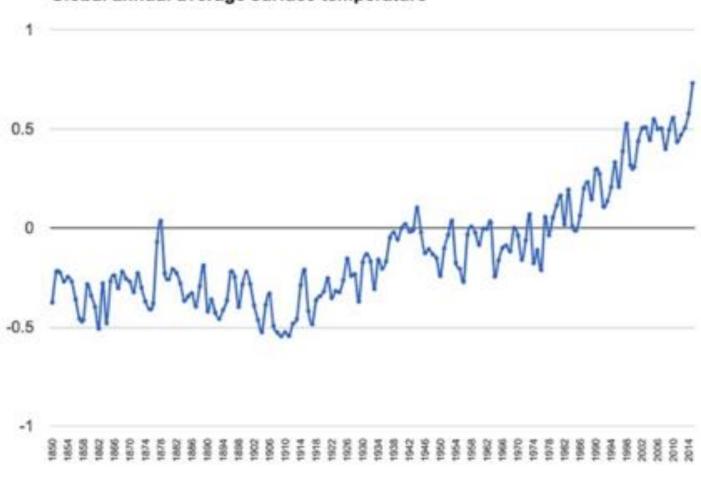
### What's Up?



# What's Up?

Annual anomaly relative to 1961-1990 (C)

#### Global annual average surface temperature



# What's Up?



### Python Scipy RV Methods

```
from scipy import stats # great package
X = stats.poisson(2.5) # X ~ Poi(λ = 2.5)
print(X.pmf(2)) # P(X = 2)
```

Function	Description
X.pmf(k)	P(X=k)
X.cdf(k)	$P(X \leq k)$
<pre>X.entropy()</pre>	(Differential) entropy of $X$
X.mean()	$\mathrm{E}[X]$
X.var()	Var(X)
X,std()	Std(X)

# **Next Time**

#### Discrete Distributions

#### Bernoulli:

indicator of coin flip X ~ Ber(p)

#### **Binomial**:

# successes in n coin flips X ~ Bin(n, p)

#### Poisson:

# successes in n coin flips X ~ Poi(λ)

#### Geometric:

# coin flips until success X ~ Geo(p)

#### **Negative Binomial:**

# trials until r successes X ~ NegBin(r, p)

#### Zipf:

- The popularity rank of a random word, from a natural language
- X ~ Zipf(s)

#### The Poisson Common Path

